

# A Performance Evaluation Protocol for Graphics Recognition Systems

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**Abstract.** This paper defines a computational protocol for evaluating the performance of raster to vector conversion systems. The graphical entities handled by this protocol are continuous and dashed lines, arcs, and circles, and text regions. The protocol allows matches of the type one-to-one, one-to-many, and many-to-one between the ground truth and the recognition results.

## 1 Introduction

Systems which convert existing paper-based drawings into electronic format are in demand and a few have been developed. However, the performance of the prototypes and commercial systems is either unknown, or only reported in a limited way by the system developers. An evaluation for these systems, or their sub-systems, would contribute to the advancement of the field. Responding to this need, a dashed-line detection competition for developers of dashed-line detection algorithms was proposed and took place during the first International Workshop on Graphics Recognition at Penn State University, in 1995. A benchmark[1] was developed and used in that competition. That benchmark includes a performance evaluator and a software tool that automatically generates dashed-line test images and the corresponding groundtruth.

In this paper, we extend that protocol to evaluate the performance of graphics recognition systems on images that contain straight lines (solid or dashed), circles (solid or dashed), partial arcs of circles (solid or dashed), and text blocks. (Engineering drawings primarily use a combination of these geometric elements. Therefore, despite being restricted to these simple entity types, the evaluation protocol is applicable to a wide variety of drawings. Upgrading the evaluator to handle other types of entities is straight forward. To do this, one needs to provide the evaluator the parameters of the entity or entities and the performance evaluation criteria.) To evaluate a given recognition system, the system is tested on a set of pre-selected test images. (The groundtruth for the test images must be reliable.) The results of the recognition system are matched, using the criteria defined in this protocol, with the corresponding groundtruth of the test

images. The matching results are the numbers of one-to-one matches, one-to-many matches, many-to-one matches, as well as the numbers of false-alarms and misses. Performance measurements for the recognition system can be formulated using a linear combination of some or all of the matching results.

Our evaluator is designed to be used by recognition system researchers and developers for testing and enhancing their recognition algorithms. The evaluator allows the users to select ‘text-only’, ‘graphics-only’, or ‘all’ option for their systems performance evaluations. The ‘text-only’ evaluation option is designed for recognition systems that detect only the text blocks in the input image. The ‘graphics-only’ evaluation option is designed for recognition systems that detect only the graphical entities. The ‘all’ evaluation option is designed for systems that can detect both graphics and text blocks.

This paper is organized as follows: In Sect. 2, we give a brief review of some related work. In Sect. 3 we specify the parameters for the entities. The protocols for performance evaluation and entity matching are given in Sect. 4. In Sect. 5, we present the matching criteria for each pair of valid combinations. The precise definitions of the line-to-line matching functions are given in Appendix A.

## 2 Previous Work

Performance evaluation and benchmarking have been gaining acceptance in all areas of computer vision. An overview of this area is available at [2]. Performance evaluation of graphics recognition is still a very young field; objective and quantitative methods for evaluation of graphics recognition have been proposed very recently [1, 3, 4, 5]. Kong et al. [1] propose a quantitative method for evaluating the recognition of dashed lines. Hori and Doermann [3] propose a quantitative performance measurement methodology for task-specific raster to vector conversion. Wenyin and Dori [4] present a protocol for evaluating the recognition of straight and circular lines. All of these methods are limited in their applicability.

Kong et al. [1] use angle, distance, relative overlap, and offset between line segments for evaluating line matches and for detecting line styles. They use several arbitrary and rigid thresholds. They do not allow for fragmentation of detected lines.

Hori and Doermann [3] instantiate and extend Haralick’s framework for performance characterization in image analysis [6], in an application-dependent manner, for measuring the performance of raster to vector conversion algorithms. The ‘applications’ addressed in the work are thinning, medial line finding, and line fitting – all low level techniques that do not completely constitute vectorization. It is hard to extend the work to evaluate a complete vectorization system. Hori and Doermann’s protocol does not distinguish between detection rate and false alarm rate. It does not include an overall evaluation metric. It does not allow for fragmentation of detected lines.

Wenyin and Dori [4] propose performance evaluation indices for straight and circular line detection. Detection and false alarm rates are defined at both the pixel level and the vector level. Pixel level performance indices (measures of shape preservation) are not appropriate when dealing with real images that

contain severe distortion introduced by warping and other defects in the hard copy drawing and by the scanning/imaging system. Attempts to obtain a high pixel recovery index would unnecessarily require the detected vectors to be true to the distorted shape of the imaged lines, thereby making the detected lines fragmented. Wenyin and Dori weight all true positives and false positives by their respective lengths. This is inappropriate if the goal of the evaluation is to measure the cost of post-processing operations that are necessary to correct the mistakes of vectorization. The time for manual post processing does not depend significantly on the length of a true positive, a missed entity, or a false positive. Time taken for adding or deleting a line in a CAD tool does not depend significantly on the length of the line.

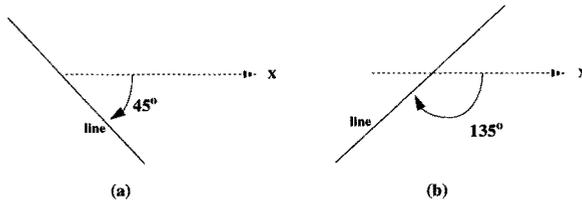
Neither of the above methods addresses the extraction or separation of text from graphics. It is not possible to evaluate graphics recognition systems on realistic drawings without accounting for text in the drawings. Wenyin and Dori [5] propose a protocol for evaluating text-graphics separation. In this protocol, the quality of the recognized text boxes is measured using  $Q_b$ , the basic quality, and  $Q_{fr}$ , the fragmentation quality. The protocol does not explicitly penalize overlapping text boxes; they are penalized in an indirect way.  $Q_{fr}$  implicitly penalizes overlap among recognized text boxes. For  $N$  recognized text boxes that are identical and have a 100% overlap with a text box in the groundtruth,  $Q_{fr}$  would be  $1/\sqrt{N}$ . The theoretical basis for penalizing overlapping text box recognition by  $1/\sqrt{N}$  is not stated. Moreover, the protocol of [5] does not allow one to accept one of the  $N$  identical text boxes as a good match and to label others as false alarms. Therefore, this penalty term cannot be used to measure post-processing/editing cost.

### 3 Entity Definition

#### 3.1 Entity Specification

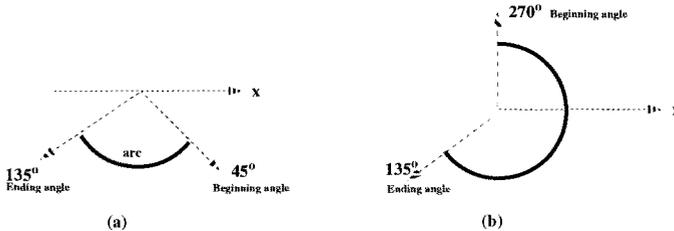
Currently our evaluator handles seven types of entities: solid and dashed lines, solid and dashed arcs, solid and dashed circles, and text areas. The specification of the parameters for these seven types are given below.

- Solid or dashed line type: For a solid or dashed line segment, the parameters are: the entity type indicator (a solid line or a dashed line), the x- and y-coordinates (this is equivalent to the c- and r-coordinate system of the image coordinate system given in Appendix A) of the two end points (no special ordering for the two points) and the orientation of the line. The orientation we use here is in degree, clock-wise with respect to x-axis. (See Fig. 1.)
- Solid or dashed circle type: For a solid or dashed circle, the parameters are: the entity type indicator (a solid or a dashed circle), the x- and y- coordinates of the center, the radius, and the thickness of the circle arc.
- Solid or dashed arc (partial circle) type: For a solid or dashed arc, the parameters are: the entity type indicator (a solid or a dashed arc), the x- and y- coordinates of the center, the radius, the beginning and the ending angles



**Fig. 1.** The orientation of the line in (a) is 45 degrees, in (b) is 135 degrees.

(in a clock-wise order, in degrees) of the arc. The orientation of the beginning and ending angles for arcs are also in degrees, clock-wise with respect to x-axis. In some cases, the beginning angles can be larger than the ending angles. (See Fig. 2.)



**Fig. 2.** The beginning and ending angles of two arcs. In (b), the beginning angle of the arc is larger than the ending angle of the arc.

- Text area type: A text area is represented by a rectangular box and its orientation. The parameters are: the entity type indicator (a text area), the x- and y- coordinates of any two opposite corners of the rectangle, and the orientation of the longest side of the rectangle.

### 3.2 Valid Entity Type Combinations

To speed up the matching score computation, we compute only those pairs having potential matches (e.g., line with line, etc.). The matching score for all incomparable combinations are set to zero. The following is the list of the valid entity type combinations. Others combinations are considered as incomparable.

- Solid-line with solid-line
- Solid-circle with solid-circle
- Solid-arc with solid-arc
- Dashed-line with dashed-line
- Dashed-circle with dashed-circle
- Dashed-arc with dashed-arc
- Solid-line with solid-arc

- Dashed-line with dashed-arc
- Solid-arc with solid-circle
- Dashed-arc with dashed-circle
- Text-area with text-area

## 4 Performance Evaluation Protocol

### 4.1 Evaluation Overview

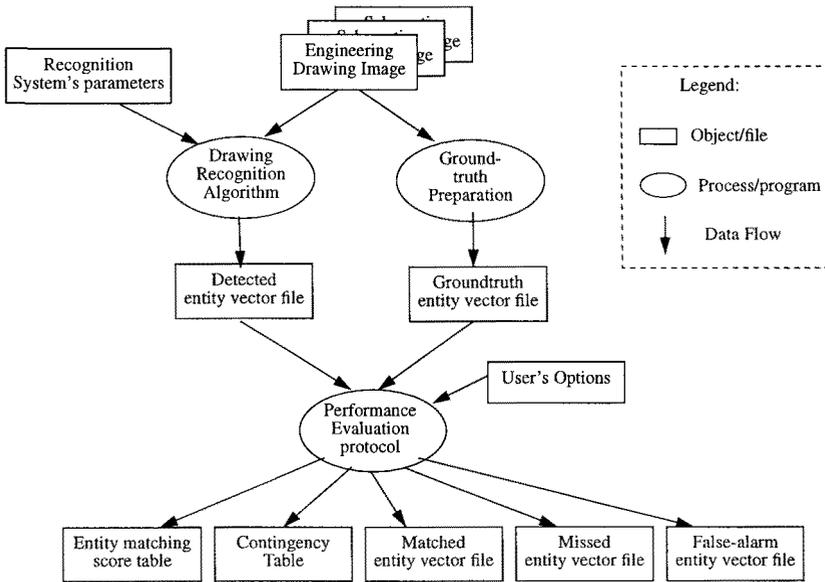


Fig. 3. The object-process diagram of our evaluator

Inputs to the evaluator are entities (in ASCII vector format) of the recognition algorithm's output and the corresponding *groundtruth*. (Figure 3 shows an object-process diagram of our evaluator.) Since there are seven types of entities (solid and dashed lines, solid and dashed arcs, solid and dashed circles, and text areas) that are allowed, the evaluation protocol and the matching criteria are designed differently for each of the combinations. The matching scores for each pair is computed according to the pair's entity type combination, using the matching criteria defined for this combination. (These criteria are defined in Sect. 5.) A match-score table is produced from the matching score computation.

From the computed match score table, we search for all the one-to-one matches, resolving the problems of the one-to-many matches and the many-to-one matches, as well as, those false-alarms and misses. The matching results are the numbers of one-to-one matches, one-to-many matches, many-to-one matches,

as well as the numbers of false-alarms and misses. Performance measurements for the recognition system can be formulated, using a linear combination of some or all of the matching results, which when weighted by application specific weights can be summed to produce an overall score relevant to the application.

## 4.2 Evaluation Protocol

The performance (accuracy) of a detection algorithm can be measured by counting the number of matches between the entities detected by the algorithm and the entities in the groundtruth, and the numbers of misses and false alarms. We consider a perfect result of a detection algorithm, if each and every one of the entities in the detected list matches one and only one entity of the same type in the groundtruth list and vice versa. The following is the protocol of this computation.

- Step 1: Obtain the detected entities and the entities' parameters and form a detected entities list (D-list) for the entities and the parameters.
- Step 2: Obtain the groundtruth entities and the entities' parameters and form a groundtruth entities list (G-list) for the entities and the parameters.
- Step 3: Compute the matching score table. (See Sect. 4.3).
- Step 4: Compute the one-to-one matches, resolving the problems of several matches, either from the detections or from the groundtruth. (See Sect. 4.4).
- Step 5: Compute the one-to-many and many-to-one partial matches.
- Step 6: Compute the false-alarms and the misses. (See Sect. 4.6).

## 4.3 Matching Score Table Computation

The matching score table is computed as follows. We compute the matching scores (ranging from 0 to 1, 1 being a perfect match) for each pair (with a valid combination) of entities, one from D-list (detected entities) and one from G-list (groundtruth entities), using the matching criteria defined (in Sect. 5) for the pair's combination. The matching score for all invalid combinations are set to zero. A two-dimensional data structure, the *match-score table*, is used to store the results of this computation. A higher matching score indicates a higher degree of match between the corresponding pair of entities. Figure 4 illustrates such a table. Blanks are read as zeros.

Note that entries in each row  $i$  of the match-score table represent the matching results from the  $i$ -th entity in the D-list to all entities in the G-list. Within a given row, a single entry having a high score value indicates a potential good match of the pair corresponding to that entry.

## 4.4 Computing One-to-One Matches

The protocol for computing the entity one-to-one matches is as follows.

**Match-Score Table**

|    | g1  | g2  | g3  | g4  | g5  | g6 | g7  | g8 | g9  | g10 |
|----|-----|-----|-----|-----|-----|----|-----|----|-----|-----|
| d1 |     |     |     |     |     |    | .85 |    | .14 |     |
| d2 |     |     |     | 1.0 |     |    |     |    |     |     |
| d3 |     |     | .1  |     |     |    | .9  |    | .1  |     |
| d4 |     |     |     |     | .95 |    | .9  |    |     |     |
| d5 | .25 | .3  | .86 |     |     |    |     | .3 | .88 |     |
| d6 |     | 1.0 |     |     |     |    |     |    |     |     |
| d7 |     | .06 | .91 |     |     |    |     |    | .93 |     |
| d8 |     | .91 |     |     |     |    |     |    |     |     |

**Fig. 4.** An example of a match-score table

Step 1: We compute a two-dimensional match-count table from the computed match-score table. The entry  $\text{match-count}(i, j)$  is set to 1 if the  $\text{match-score}(i, j)$  is greater or equal to *upper-threshold*, otherwise, it is set to zero. Currently, the *upper-threshold* is set to .85. However, we allow users to set their own threshold. Figure 5 illustrates the match-count table (on the right) which is computed from the match-score table on the left.

**Match-Score Table**

|    | g1  | g2  | g3  | g4  | g5  | g6 | g7  | g8 | g9  | g10 |
|----|-----|-----|-----|-----|-----|----|-----|----|-----|-----|
| d1 |     |     |     |     |     |    | .85 |    | .14 |     |
| d2 |     |     |     | 1.0 |     |    |     |    |     |     |
| d3 |     |     | .1  |     |     |    | .9  |    | .1  |     |
| d4 |     |     |     |     | .95 |    | .9  |    |     |     |
| d5 | .25 | .3  | .86 |     |     |    |     | .3 | .88 |     |
| d6 |     | 1.0 |     |     |     |    |     |    |     |     |
| d7 |     | .06 | .91 |     |     |    |     |    | .93 |     |
| d8 |     | .91 |     |     |     |    |     |    |     |     |

→

**Match-count Table**

|    | g1 | g2 | g3 | g4 | g5 | g6 | g7 | g8 | g9 | g10 |
|----|----|----|----|----|----|----|----|----|----|-----|
| d1 |    |    |    |    |    |    | 1  |    |    |     |
| d2 |    |    |    | 1  |    |    |    |    |    |     |
| d3 |    |    |    |    |    |    | 1  |    |    |     |
| d4 |    |    |    |    | 1  |    | 1  |    |    |     |
| d5 |    |    | 1  |    |    |    |    |    | 1  |     |
| d6 |    | 1  |    |    |    |    |    |    |    |     |
| d7 |    |    | 1  |    |    |    |    |    |    | 1   |
| d8 |    | 1  |    |    |    |    |    |    |    |     |

**Fig. 5.** An example of a match-count table (on the right)

Step 2: Two projection profiles (D-profile and G-profile) are computed from the match-count table, the result of Step 1.

The entry  $D(i)$  is computed as the sum of the matches in the  $i$ -th row of the match-count table. Likewise, the entry  $G(j)$  is computed as the sum of the matches in the  $j$ -th column of the match-count table. Figure 6 illustrates the D-profile and the G-profile computed from the match-count table.

The interpretation of these two profiles can be as follows:

- One-to-one matches: An  $i$ -th D entity has a one-to-one match with a  $j$ -th G entity, if
  - $D(i)$  is a one,
  - the  $\text{match-count}(i, j)$  is one,
  - and  $G(j)$  is one.

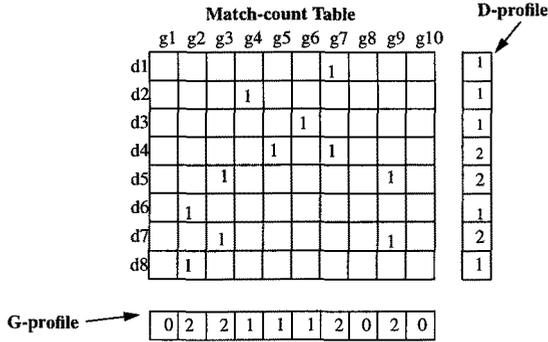


Fig. 6. The two projection profiles for the match-count table of Fig. 5

(If the detection algorithm produces a perfect result, all entries in the D-profile and the G-profile will be one.)

- Many-to-one conflicts: An entry in G-profile, say  $G(j)$ , is greater than one. That is, there are multiple D entities matching with the  $j$ -th entity in G-list.
- One-to-many conflicts: An entry in D-profile, say  $D(i)$ , is greater than one. That is, the  $i$ -th D entity matches two or more entities in the G-list.
- False-alarms: A zero entry in the D-profile indicates that no strong match is found from this D entity to any of the entities in the G-list.
- misses: A zero entry in the G-profile indicates that no strong match is found from this G entity to any of the entities in the D-list.

Step 3: Compute the one-to-one match list.

For each  $D(i)$  that is equal to one, we attempt to locate the pair  $(i, j)$  such that both  $G(j)$  and  $\text{match-count}(i, j)$  are one, a one-to-one match (d2 and d3 in Fig. 7). We put the pair,  $(i, j)$ , in the one-to-one match list, and set  $D(i)$  and  $G(j)$  to -1 (block the two entities from further consideration). Figure 7 illustrates the result of one-to-one matching. There are two such pairs.

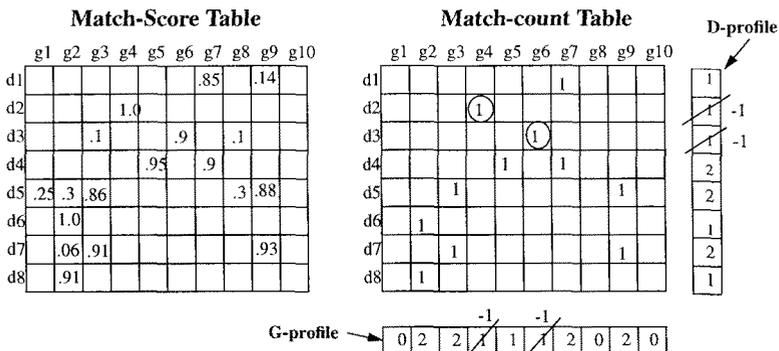


Fig. 7. One-to-one matching (the resulting one-to-one matches are circled)

Step 4: Resolving the many-to-one conflicts.

For each  $D(i)$  that is equal to one (but did not produce a one-to-one match in Step 3), we locate the pair  $(i, j)$  such that  $\text{match-count}(i, j)$  is one, but  $G(j)$  is greater than one. (There are three such pairs in Fig. 7):  $(d1, g7)$ ,  $(d6, g2)$ , and  $(d8, g2)$ .)

Without loss of generality, let  $D(i)$  and  $D(k)$  be one and let  $G(j)$  be two (meaning that there are two  $D$  entities, say  $i$ -th and  $k$ -th, matching with the  $j$ -th entity in  $G$ -list, such as  $(d6, g2)$ , and  $(d8, g2)$  in Fig. 7),

Case 1: We select the pair  $(i, j)$  if  $\text{match-score}(i, j) \geq \text{match-score}(k, j)$ .

And

- we put the pair  $(i, j)$  in the one-to-one match list;
- we set  $D(i)$  and  $G(j)$  to -1, and
- we decrease  $D(k)$  by one.

For example, if  $i = 6$ , we would select the pair  $(d6, g2)$  over  $(d8, g2)$  in Fig. 7.

Case 2: We select the pair  $(k, j)$  if  $\text{match-score}(k, j) > \text{match-score}(i, j)$  and  $D(k)$  is a one. And,

- we put the pair  $(k, j)$  in the one-to-one match list;
- we set  $D(k)$  and  $G(j)$  to -1, and
- we decrease  $D(i)$  by one.

The example, if  $i = 8$ , we still select the pair  $(d6, g2)$  over  $(d8, g2)$  in Fig. 7.

Case 3:  $\text{match-score}(i, j) < \text{match-score}(k, j)$  and  $D(k)$  is greater than one,  $(d1$  and  $d4$  with  $g7$ , in Fig. 7).

In this case, we would not select the pair  $(k, j)$  if there is a column  $t$ , in row  $k$  such that  $\text{match-score}(k, t) > \text{match-score}(k, j)$ . In this case, we would select the pair  $(i, j)$  instead. And we handle this case as in Case 1.

The example in Fig. 7,  $d1$  and  $d4$  match  $g7$ , a many-to-one conflict. The pair  $(d1, g7)$  is selected instead of  $(d4, g7)$  since the  $\text{match-score}(d4, g5)$  has a higher score than the pair  $(d4, g7)$ .

A similar treatment is done if  $G(j)$  is greater than two. This step is repeated until no more  $D(i)$  is equal to one.

Step 5: Resolving the one-to-many conflicts.

For each  $D(i)$  that is a two, let  $j$  and  $k$  be the two entities in  $G$ -list that match with the  $i$ -th entity in the  $D$ -list. If  $\text{match-score}(i, j) \geq \text{match-score}(i, k)$ , we put the pair  $(i, j)$  in the one-to-one match list and set  $D(i)$  and  $G(j)$  to -1, and decrease  $G(k)$  by one. Otherwise, we put the pair  $(i, k)$  in the one-to-one match list and set  $D(i)$  and  $G(k)$  to -1, and decrease  $G(j)$  by one. A similar treatment is done if  $G(j)$  is greater than two. This step is repeated until no more  $D(i)$  is two or greater.

#### 4.5 Computing Partial Matches: One-to-Many and Many-to-One

At the end of the one-to-one entity matching (Sect. 4.4), the value of each  $D(i)$  indicates whether the  $i$ -th entity has a one-to-one match. In particular, if a  $D(i)$  is

a -1, it indicates that the  $i$ -th entity in  $D$ -file has a one-to-one match, otherwise, it indicates that it does not have a one-to-one match. (The same idea for the entities in  $G$ -file.) One could, for example, consider each  $D(i) \geq 0$  a false alarm and each  $G(j) \geq 0$  a miss detection.

However, it may be the case that a detected entity which does not have a one-to-one match may in fact match with a group of two or more groundtruth entities. For example, a detection algorithm may have located a text bounding box on the input image that includes several text lines, while those text lines are given one text bounding box each in the groundtruth file.

Therefore, to give partial credits to the detection algorithms for finding one-to-many partial matches, we do as follows.

For each  $D(i) \geq 0$  in the  $D$ -profile, we collect a list of all entities  $j$  in  $G$ -profile, such that  $G(j)$  is also  $\geq 0$  (also did not have a match by any  $D$  entity) and the score in the  $match-score(i, j)$  is greater than the *lower-threshold* (so that we would not include any noise). Currently, the *lower-threshold* is set to .05. However, we allow users to set their own threshold. If the sum of the scores for the entities in the collected list is greater than the *upper-threshold*, we consider the  $i$ -th  $D$  entity having a one-to-many partial match to those  $j$  entities in the collected list. And we set  $D(i)$  and all those  $G(j)$  to -1.

A similar protocol is applied to find the many-to-one matches (many detected entities matching with one groundtruth entity).

## 4.6 False-alarms and Miss-Detections

Finally, the false-alarms are those  $i$  entities having their  $D[i]$  values  $\geq 0$ , and the miss-detections are those  $j$  entities having their  $G[j]$  values  $\geq 0$ .

# 5 Matching Criteria

## 5.1 Line-Line Matching Protocol and Criteria

This protocol is for both solid lines and dashed lines. Let  $d$  be a line entity in the  $D$ -entity-list and  $g$  be a line entity in the  $G$ -entity-list. Let  $match-score(i, j)$  be the corresponding entry of  $d$  and  $g$ . To mark the entry  $(i, j)$ , we compute the followings:

- Step 1: If  $d$  and  $g$  have the same endpoints (a perfect match), we set  $match-score(i, j)$  to 1 and skip the following steps.
- Step 2: We compute the angle between  $d$  and  $g$  as the included angle between these two line segments (see Fig. 8). If the *angle* between  $d$  and  $g$  is less than 5 degrees,  $angle(d, g) \leq 5$ , we continue to the next step, otherwise the  $match-score(i, j)$  is set to zero.
- Step 3: We compute the *distance*,  $llDist(d, g)$ , between  $d$  and  $g$ . The  $llDist(d, g)$  is computed as the average of the orthogonal distance from the midpoint of  $d$  to  $g$  and the orthogonal distance from the midpoint of  $g$  to  $d$  (see Fig. 9). If  $llDist(d, g) \leq \theta_u$ , we continue to the next step, otherwise the *entity-match-table* is set to zero.

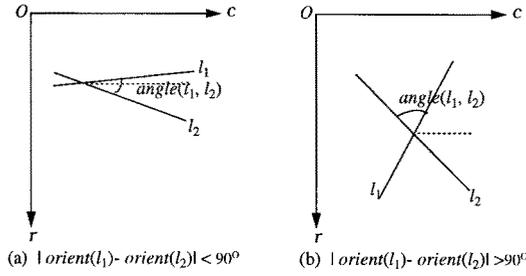


Fig. 8. Angles of two line segments

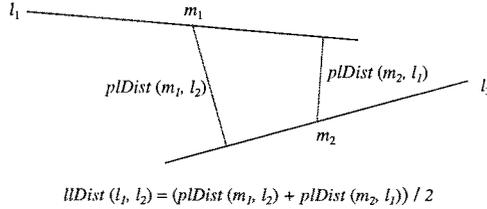


Fig. 9. Line-line distance of two line segments

Step 4: Next, we compute the *relative overlap* function,  $overlap(d, g, \alpha_g)$ , of  $d$  and  $g$  with respect to the orientation of  $g$ . If  $overlap(d, g, \alpha_g)$  is at least 20% of both  $d$  and  $g$ , we set  $match-score(i, j)$  to zero. Otherwise, we compute the relative overlap.

The relative overlap of two line segments  $l_1$  and  $l_2$  is defined as the ratio between the overlap function  $overlap(l_1, l_2, \alpha)$  and the length of the longer segment:  $RelativeOverlap(l_1, l_2, \alpha) = \frac{overlap(l_1, l_2, \alpha)}{\max(length(l_1), length(l_2))}$

Step 5: The  $match-score(i, j)$  is computed as:

$$RelativeOverlap(l_1, l_2, \alpha) - \frac{angle(d, g)}{180} - \frac{llDist(d, g)}{\theta_{11}}$$

For a precise definition of these functions see the Appendix A. The threshold values used here were determined heuristically based on the dimension of the entities. However, we allow these threshold values to be set by the user.

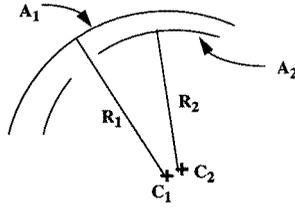
### 5.2 Arc-Arc Matching Protocol and Criteria

This protocol is for both solid arcs and dashed arcs.

Let  $A_1$  be an arc entity in the  $D$ -entity-list and  $A_2$  be an arc entity in the  $G$ -entity-list. Let  $C_1$  and  $C_2$  be the centers of  $A_1$  and  $A_2$ , and let  $R_1$  and  $R_2$  be the two radii (see Fig. 10).

Let  $match-score(i, j)$  be the entry that stores the matching result for this pair. The protocol for computing  $match-score(i, j)$  is as follows:

Step 1: If the two arcs,  $A_1$  and  $A_2$ , are identical (with the same centers, same radii, same beginning and end angles), we set  $match-score(i, j)$  to 1 and skip the following steps.



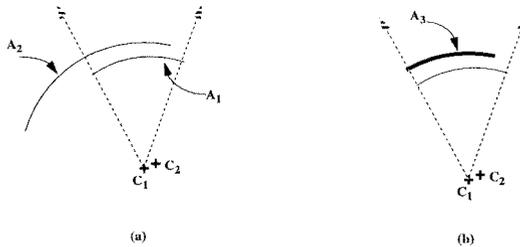
**Fig. 10.** An example of arc-arc entity pair

Step 2: We compute the point-to-point distance,  $ppDist(C_1, C_2)$ , between the two centers  $C_1$  and  $C_2$ . If this distance is greater than  $\theta_{CC}$ , we set match-score(i, j) to zero, and skip the following steps.

Step 3: We compute the absolute difference between the two radii,  $RadiusDist(R_1, R_2) = |R_1 - R_2|$ . If this distance is greater than  $\theta_{RR}$ , we set match-score(i, j) to zero, and skip the following steps.

Step 4: We compute the ratio of these two radii,  $RadiusRatio(R_1, R_2) = \frac{\min(R_1, R_2)}{\max(R_1, R_2)}$ . If this ratio is smaller than 85 percent, we set match-score(i, j) to zero, and skip the following steps.

Step 5: We project  $A_1$  onto  $A_2$ , take the portion of  $A_2$  that is within this projection, and call it  $A_3$ . The projection is defined as follows. We construct two lines,  $l_1$  and  $l_2$ , from the center of  $A_1$  to the two endpoints of  $A_1$  to infinity. The wedge between these two lines (also within the beginning and the ending angles of  $A_1$ ) is the projection of  $A_1$ . (See Fig. 11).

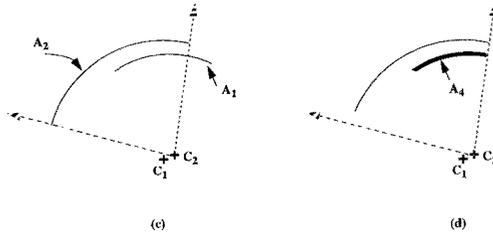


**Fig. 11.** (a) The projection of  $A_1$  onto  $A_2$ , (b)  $A_3$ , the portion of  $A_2$  within the projection

Similarly, we project  $A_2$  onto  $A_1$  and take the portion of  $A_1$  that is within this projection, call it  $A_4$  (see part (d) in Fig. 12).

Step 6: If either  $A_3$  or  $A_4$  exists, we set the match-score(i, j) to zero and skip the rest. In the case that Both  $A_3$  and  $A_4$  exist, we define a line segment  $L_1$  from the two ends of  $A_3$  and another line segment  $L_2$  from the two ends of  $A_4$ .

Step 7: Taking  $L_1$  and  $L_2$ , we apply the line-line matching criteria to this pair.



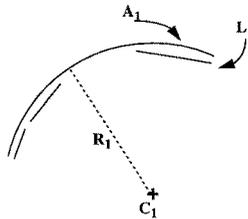
**Fig. 12.** (c) The projection of  $A_2$  onto  $A_1$ , (d)  $A_4$ , the portion of  $A_1$  within the projection

The match-score( $i, j$ ) for this pair is the result of the line-line matching result times  $F$ , an adjustment. Without loss of generality, let  $L_1$  be shorter than  $L_2$  and the distance between the two ends of  $A_2$  be greater than that of  $A_1$ . Then  $F$  is computed as the ratio of the length of  $L_1$  and the distance between the two ends of  $A_2$ .

### 5.3 Arc-Line Matching Protocol and Criteria

This protocol is for solid-arc with solid-line pairs and dashed-arc with dashed-line pairs.

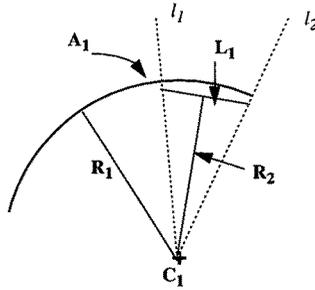
Let  $A_1$  be an arc entity in the *D-entity-list*, and let  $C_1$  be the center and  $R_1$  be the radius of  $A_1$ . Let  $L_1$  be a line entity in the *G-entity-list* (see Fig. 13).



**Fig. 13.** An arc-line entity pair

Let match-score( $i, j$ ) be the entry that stores the matching result for this pair. The protocol for computing match-score( $i, j$ ) is as follows:

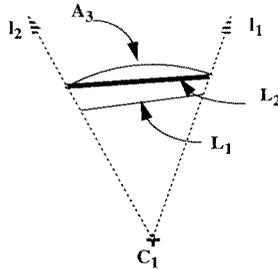
Step 1: We construct two lines,  $l_1$  and  $l_2$ , from the center of  $A_1$ , to the two endpoints of  $L_2$ . We also construct a new line segment,  $R_2$ , from the center of  $A_1$  to the midpoint of  $L_2$ . In the sense, we are constructing an artificial arc,  $A_2$ , taking  $C_1$  as its center,  $R_2$  as its radius, and the orientation of  $l_1$  as its beginning angle, and the orientation of  $l_2$  as its ending angles (see Fig. 14).



**Fig. 14.** Construction of the new triangle and the new line segment

Step 2: We compute the absolute difference between  $R_1$  and  $R_2$ ,  $|R_1 - R_2|$ . If this distance is greater than  $\theta_{RR}$ , we set  $\text{match-score}(i, j)$  to zero, and skip the following steps.

Step 3: We project the artificial arc,  $A_2$ , onto  $A_1$ , take the portion of  $A_1$  between  $l_1$  and  $l_2$ , and call it  $A_3$ . If  $A_3$  does not exist, we set  $\text{match-score}(i, j)$  to zero and skip the rest. In the case that  $A_3$  exists, we construct a new line segment,  $L_1$ , from the new arc  $A_3$  (see Fig. 15).



**Fig. 15.** The new line segment,  $L_1$ , which is constructed from  $A_3$

Step 4: Taking  $L_1$  and  $L_2$ , we apply the line-line matching criteria to this pair.

The match-score is set to the result of the line-line matching result times  $F$ , an adjustment. Without loss of generality, let  $L_1$  be shorter than  $L_2$  and the distance between the two ends of  $A_1$  be greater than that of  $A_2$ . Then  $F$  is computed as the ratio of the length of  $L_1$  and the distance between the two ends of  $A_1$ . If either  $L_1$  or  $L_2$  exist, we set the  $\text{match-score}(i, j)$  to zero.

#### 5.4 Arc-Circle Matching Protocol and Criteria

This protocol is for solid-circle with solid-arc pairs and dashed-circle with dashed-arc pairs.

Let  $A_1$  be an arc entity in the *D-entity-list* and  $C_{ir}$  be a circle entity in the *G-entity-list*. Let  $C_1$  and  $C_2$  be the centers of  $A_1$  and  $C_{ir}$ , and let  $R_1$  and  $R_2$  be the two radii.

Let  $\text{match-score}(i, j)$  be the entry that stores the matching result for this pair of entities. The protocol for computing  $\text{match-score}(i, j)$  is as follows:

Step 1: We project  $A_1$  onto  $Cir_2$ , take the portion of the  $Cir_2$  that is within this projection, and call it  $A_2$ . In the sense, we are constructing an artificial arc,  $A_2$ , taking  $C_2$  as its center,  $R_2$  as its radius, and the orientations of the two sides of the projection as its beginning and ending angles.

Step 2: We construct a line segment,  $L_1$ , from the two ends of  $A_1$  and another line segment,  $L_2$ , from the two ends of  $A_2$ . Taking  $L_1$  and  $L_2$ , we apply the line-line matching criteria to this pair. The match-score is set to the result of the line-line matching result times  $F$ , an adjustment.  $F$  is computed as the inside angle of  $A_1$  divided by 360.

## 5.5 Circle-Circle Matching Protocol and Criteria

This protocol is for both solid and dashed circles.

Let  $Cir_1$  be a circle entity in the  $D$ -entity-list and  $Cir_2$  be a circle entity in the  $G$ -entity-list. Let  $C_1$  and  $C_2$  be the two centers and let  $R_1$  and  $R_2$  be the two radii (see Fig. 16).

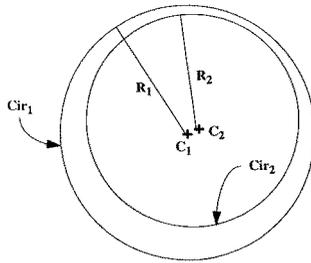


Fig. 16. A circle-circle entity pair

Let  $\text{match-score}(i, j)$  be the entry that stores the matching result for this pair of entities. The protocol for computing  $\text{match-score}(i, j)$  is as follows:

Step 1: If the two circles,  $Cir_1$  and  $Cir_2$ , are identical (with the same centers and the same radii), we set  $\text{match-score}(i, j)$  to 1 and skip the following steps.

Step 2: We compute the point-to-point distance,  $ppDist(C_1, C_2)$ , between the two centers  $C_1$  and  $C_2$ . If this distance is greater than  $\theta_{CC}$ , we mark the entry as a non-match and skip the following steps.

Step 3: We compute the absolute difference between the two radii,  $RadiusDist(R_1, R_2) = |R_1 - R_2|$ . If this distance is greater than  $\theta_{RR}$ , we mark the entry as a non-match and skip the following steps.

Step 4: If both  $ppDist(C_1, C_2)$  and  $RadiusDist(R_1, R_2)$  are equal to zero (two identical circles), we set  $match-score(i, j)$  to 1 and skip the following steps.

Step 5: We compute the ratio of these two radii,  $RadiusRatio(R_1, R_2) = \frac{\min(R_1, R_2)}{\max(R_1, R_2)}$ . If this ratio is smaller than the preset threshold, we mark the entry as a non-match, otherwise, the entry is marked as a match.

Step 6: The matching score for this pair is

$$RadiusRatio(R_1, R_2) - \frac{ppDist(C_1, C_2)}{\min(R_1, R_2)} - \frac{RadiusDist(R_1, R_2)}{\min(R_1, R_2)}$$

## 5.6 Text-text Matching Protocol and Criteria

A text area is represented by a rectangular box and its orientation. The rectangular box is described by the x- and y- coordinates of any two opposite corners of the rectangle, and the orientation of the longest side of the rectangle.

Let  $t_1$  be a text entity in the *D-entity-list* and let  $t_2$  be a text entity in the *G-entity-list*. Let  $P_1$  and  $P_2$  be the two opposite corners of  $t_1$  and let  $Q_1$  and  $Q_2$  be the two opposite corners of  $t_2$ .

Let  $match-score(i, j)$  be the corresponding entry of  $t_1$  and  $t_2$ . The computational protocol for the  $match-score(i, j)$  is as follows:

Step 1: If  $t_1$  and  $t_2$  are identical (a perfect match), we set  $match-score(i, j)$  to 1 and skip the following steps.

Step 2: Let  $P_{mid}$  be the midpoint of the diagonal line of  $t_1$  and let  $Q_{mid}$  be the midpoint of the diagonal line of  $t_2$ . Without loss of generality, let  $ppDist(P_1, P_2) > ppDist(Q_1, Q_2)$ . That is,  $t_1$  has longer diagonal than that of  $t_2$ . We construct a circle centered at  $P_{mid}$  with radius equal to the diagonal of  $t_1$ . If both the corner points of  $t_2$ ,  $Q_1$  and  $Q_2$ , are outside of this circle (this means that there is no overlap between  $t_1$  and  $t_2$ ), we set  $match-score(i, j)$  to zero, and skip the following steps. (The step is designed to limit the search space.)

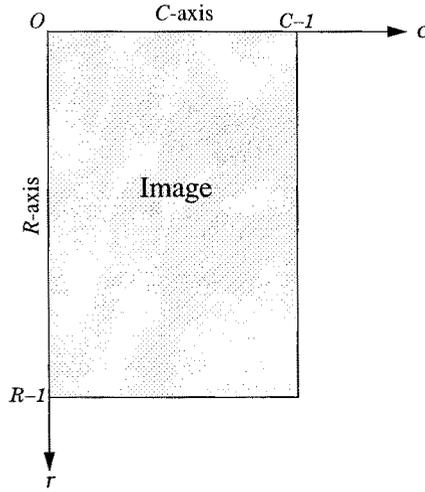
Step 3: Next, we compute the other two opposite corners of  $t_1$  and  $t_2$ , so that each of the text boxes is now represented as a rectangle. Let  $P$  and  $Q$  be the two rectangles.

Step 4: We compute the intersection of  $P$  and  $Q$ , call it,  $I$ . If  $I$  is empty, we set  $match-score(i, j)$  to zero. Otherwise, we compute the area of  $P$ ,  $Q$ , and  $I$ . And we set  $match - score(i, j) = \frac{area(I)}{\max(area(P), area(Q))}$ .

## 6 Appendix A: Line-Line Matching Criteria Functions

### Image Coordinate System

An image is given by columns and rows of pixels. In a bi-level binary image, a foreground pixel has the value 1 and a background pixel has the value 0. We use the *Column-Row* coordinate system, (*c*-coordinate, *r*-coordinate), to represent



**Fig. 17.** The column-row coordinate system for an image. The origin  $(0,0)$  is at the top-left corner pixel of the image. The image has  $R$  rows and  $C$  columns.

a pixel's position within an image. The origin of this system,  $(0,0)$ , is at the top-left corner pixel of the image (see Fig. 17).

Let  $d$  be a line entity in the  $D$ -entity-list and  $g$  be a line entity in the  $G$ -entity-list. The following functions are needed for the line-line matching criteria.

### $\alpha$ -Projection of a Line Segment

The  $\alpha$ -projection of a line  $l = (c_1, r_1, c_2, r_2)$  is the projection of  $l$  onto the given orientation  $\alpha \in (-90^\circ, 90^\circ]$ . The  $\alpha$ -projection of  $l$ ,  $proj(l, \alpha)$ , is also a line segment, its two endpoints  $(c'_1, r'_1)$  and  $(c'_2, r'_2)$  are the projections of  $(c_1, r_1)$  and  $(c_2, r_2)$  onto the orientation  $\alpha$ , respectively.

$$\begin{aligned} c'_1 &= \cos \alpha (c_1 \cos \alpha + r_1 \sin \alpha) \\ r'_1 &= \sin \alpha (c_1 \cos \alpha + r_1 \sin \alpha) \\ c'_2 &= \cos \alpha (c_2 \cos \alpha + r_2 \sin \alpha) \\ r'_2 &= \sin \alpha (c_2 \cos \alpha + r_2 \sin \alpha) \end{aligned}$$

If  $|\alpha - orient(l)| \leq 90^\circ$ ,  $proj(l, \alpha)$  is given by  $(c'_1, r'_1, c'_2, r'_2)$ ; otherwise, it is given by  $(c'_2, r'_2, c'_1, r'_1)$ .

$$proj(l, \alpha) = \begin{cases} (c'_1, r'_1, c'_2, r'_2) & \text{if } |\alpha - orient(l)| \leq 90^\circ \\ (c'_2, r'_2, c'_1, r'_1) & \text{otherwise} \end{cases}$$

## Overlap of Two Line Segments – A Relationship Function

The  $\alpha$ -overlap of two line segment  $l_1$  and  $l_2$ ,  $overlap(l_1, l_2, \alpha)$ , is a relationship function of  $l_1$  and  $l_2$  with respect to a given orientation  $\alpha$ .

Suppose  $T_a$  and  $T_d$  are two given thresholds, which are determined by application and user.  $T_a$  is a threshold for the angle of two line segments,  $T_d$  is a threshold for the line-line distance. If  $angle(l_1, l_2)$  is not greater than  $T_a$ , and  $lDist(l_1, l_2)$  is not greater than  $T_d$ , we say that  $l_1$  and  $l_2$  are sufficiently close.

The  $\alpha$ -overlap of  $l_1$  and  $l_2$  is defined as the length of the common part of their  $\alpha$ -projections if  $l_1$  and  $l_2$  are sufficiently close, and is defined as 0 otherwise. The function  $overlap(l_1, l_2, \alpha)$  is

$$overlap(l_1, l_2, \alpha) = \begin{cases} length(proj(l_1, \alpha) \cap proj(l_2, \alpha)) & \text{if } angle(l_1, l_2) \leq T_a, \text{ and } lDist(l_1, l_2) \leq T_d \\ 0 & \text{otherwise} \end{cases}$$

## Relative Overlap

The relative overlap of two line segments  $l_1$  and  $l_2$  is defined as the ratio between the overlap function  $overlap(l_1, l_2, \alpha)$  and the length of the longer segment:

$$\frac{overlap(l_1, l_2, \alpha)}{\max(length(l_1), length(l_2))}$$

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